<table>
<thead>
<tr>
<th><strong>Title</strong></th>
<th>Scheduling Methods Based on Data Division for Continuous Media Data Broadcast</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Author(s)</strong></td>
<td>Yoshihisa, Tomoki; Tsukamoto, Masahiko; 西尾, 章治郎</td>
</tr>
<tr>
<td><strong>Citation</strong></td>
<td>IEEE Pacific RIM Conference on Communications, Computers, and Signal Processing – Proceedings. 2 P.927-P.930</td>
</tr>
<tr>
<td><strong>Issue Date</strong></td>
<td>2003-08</td>
</tr>
<tr>
<td><strong>Text Version</strong></td>
<td>publisher</td>
</tr>
<tr>
<td><strong>URL</strong></td>
<td><a href="http://hdl.handle.net/11094/14072">http://hdl.handle.net/11094/14072</a></td>
</tr>
<tr>
<td><strong>DOI</strong></td>
<td></td>
</tr>
<tr>
<td><strong>rights</strong></td>
<td>c2003 IEEE. Personal use of this material is permitted. However, permission to reprint/republish this material for advertising or promotional purposes or for creating new collective works for resale or redistribution to servers or lists, or to reuse any copyrighted component of this work in other works must be obtained from the IEEE.</td>
</tr>
<tr>
<td><strong>Note</strong></td>
<td></td>
</tr>
</tbody>
</table>

*Osaka University Knowledge Archive : OUKA*

https://ir.library.osaka-u.ac.jp/repo/ouka/all/
Scheduling Methods Based on Data Division for Continuous Media Data Broadcast

Tomoki Yoshihisa  
Grad. School of Information Science and Technology, Osaka University  
yoshihisa@ist.osaka-u.ac.jp

Masahiko Tsukamoto  
Grad. School of Information Science and Technology, Osaka University  
tuka@ist.osaka-u.ac.jp

Shojiro Nishio  
Grad. School of Information Science and Technology, Osaka University  
nishio@ist.osaka-u.ac.jp

Abstract—Recently, various schemes have arisen to broadcast continuous media data such as audio and video. Some of them have focused on reducing clients' waiting time under the continuity condition, i.e., to play data completely without any interruptions. These schemes usually employ multiple channels to broadcast continuous media data. However, receivers of many existing broadcast systems such as those equipped with wireless LAN or Bluetooth cannot receive data from multiple channels concurrently. In this paper, we propose and evaluate a scheduling scheme to reduce the waiting time of clients with a single channel.

I. INTRODUCTION

Recently, broadcasting continuous media data such as audio and video has become increasingly popular [1-5]. In general broadcast systems, a server broadcasts data repetitively, and clients selectively receive broadcast data which may be desired by a user. In such systems, the user orders his/her client to receive a data, then the client waits until the data is broadcast. Hence, the schedule of the broadcast data at the server side has a great effect on clients' waiting time.

Meanwhile, it is important for users to play data without any interruptions. Previous studies have aimed at reducing the waiting time of clients to start playing data under this continuity condition and this is achieved by dividing data into several segments and broadcasting these segments via multiple channels. However, with multiple channels, clients are forced to receive data from more than one channel concurrently, while the server has to control multiple channels. Hence, the broadcasting mechanism becomes more complicated than that of broadcasting data with a single channel. Furthermore, typical popular commercial wireless LAN transceivers, such as Bluetooth [6] or set-top-boxes, cannot always concurrently receive data from multiple channels.

In this paper, we propose a scheduling scheme to reduce the waiting time of clients with a single channel. In our scheme, assuming that users can play data continuously, data is divided into several segments to produce an effective schedule. By broadcasting the first segment frequently, the waiting time is reduced. In this scheme, the number of segments and the broadcast schedule affect the waiting time. We show the schedule that gives the minimum average waiting time, in which data is divided into two segments of equal size. Where data is divided into more than two segments, we propose and evaluate a heuristic scheduling method. The remainder of this paper is organized as follows. Section II explains our basic idea. We show how broadcasting precedent segments with a single channel frequently reduce the waiting time of clients. Our assumed system environment and our scheme are explained in Section III. In Section IV, for the case of dividing data into two segments, we show the schedule that gives the minimum average waiting time. Section V explains our heuristic scheduling method for the case of dividing data into more than two segments. Section VI remarks related works, and finally, we conclude the paper in Section VII.

Fig. 1. A simple repetition scheme for a continuous media data broadcast.

II. BASIC IDEA

In this section, first, we show the average waiting time of a simple repetition scheme, and following that, we outline our scheme's basic idea.

Suppose that it takes 40 seconds to broadcast a 5-minute continuous media data. Note that this case is suitable for broadcasting a 5-minute MP3 encoded music data (the data size: 5 Mbytes) on 1-Mbps satellite or Bluetooth systems. In the case of broadcasting the data repeatedly, as shown in Fig. 1 (the simple repetition scheme), a user who requests the data to his/her client should wait 20 seconds on average until the client starts playing it. Note that the data is streaming media data, i.e., the client can start playing the data immediately after receiving it.

To explain our basic idea, we divide the data into two 2.5-minute segments. It takes 20 seconds to broadcast each segment. Let $S_1$ indicate the first segment and $S_2$ the remaining segment. The server broadcasts these segments in the order of $S_1S_2S_1$ as shown in Fig. 2. Accordingly, the broadcast cycle becomes 60 seconds. Let $T_i$ ($i=1,2,\ldots$) indicate the time slot. We assume that the broadcast cycle starts at $T_1$. That is, $S_1$ is broadcast at $T_1$ and $S_2$ is broadcast at $T_2$ in sequence. The clients' average waiting time can be calculated as follows.
When requesting the data in \( T_1 \):
The client receives \( S_1 \) broadcast at \( T_2 \) and can store \( S_2 \)
broadcast at \( T_3 \) while it plays \( S_1 \). The average waiting
time in this case is 10 seconds.

When requesting the data in \( T_2 \):
The client receives \( S_1 \) broadcast at \( T_4 \) and receives \( S_2 \)
broadcast at \( T_3 \) or \( T_5 \) as soon as it finishes playing \( S_1 \).
The average waiting time in this case is 30 seconds.

When requesting the data in \( T_3 \):
The client receives \( S_1 \) broadcast at \( T_4 \) and receives \( S_2 \)
broadcast at \( T_6 \) as soon as it finishes playing \( S_1 \). The
average waiting time in this case is 10 seconds.

As a result, the average waiting time is 17 seconds. This is
15% shorter than that of the simple repetition scheme.

In this case, the average waiting time of the broadcast schedule \( S_1S_2S_1S_1S_1S_2 \) is 12.5 seconds. This is 38%
shorter than that of the simple repetition scheme.

Thus, by dividing data into several segments and the server broadcasts precedent segments frequently, the average waiting
time is reduced.

### III. Assumed System Environment and Our Scheme

We assume that the time needed to broadcast continuous media data is shorter than the playing time. Let \( D \) indicate the playing time of data and \( D' \) the time needed to broadcast it. We define the playback ratio \( a \) by the following equation:

\[
a = \frac{D}{D'}.
\]

For example, when data is encoded by MPEG Audio Layer 3 (MP3), the data size of a 5-minute audio file is approximately
5 Mbytes. Hence, it takes 40 seconds to broadcast the data on a
1-Mbps network. Hence, the playback ratio \( a \) is \( 300/40 = 7.5 \).

#### A. Our Assumed System Environment

- A client is not able to receive data from multiple channels concurrently, i.e., the client can only receive data from a single channel.
- A server broadcasts one item of continuous media data with a single channel.
- A client who starts playing data has to be able to play the data without any interruptions, i.e., the continuity condition must be satisfied.
- \( D' \) is shorter than \( D \). The playback ratio \( a (\geq 1) \) is constant.

#### B. Our Proposed Scheme

We propose a scheduling scheme, the division-based broadcasting scheme. In this scheme, an item of broadcast data is divided into \( N \) segments \( S_i (i = 1, \ldots, N) \) of equal size.

The necessary time to broadcast a segment is \( D'/N \) (Fig. 3).

The server broadcasts \( c \) segments in a broadcast cycle. Fig. 4 shows an example of the division-based broadcasting scheme for \( c = 6, N = 3 \).

One important problem relates to how we can produce a broadcast schedule that effectively reduces the waiting time. For example, if it takes 1 minute to broadcast a 5-minute item of MP3 data \( (a = 5, D = 5 \text{ min}) \) and the data is divided into two segments \( (N = 2) \), and 4 or 5 segments are broadcast in a broadcast cycle \( (c = 4, 5) \), we can suppose the broadcast schedules shown in Table I. Since \( S_1S_2S_1S_2 \) is regarded as \( c = 2 \), this broadcast schedule is not shown in the table. The minimum average waiting time of Table I is 21 minutes and the broadcast schedule is \( S_1S_2S_1S_2 \); thus, for any given \( a \), \( c \), and \( N \), the average waiting time is calculated. However, since these values can be given arbitrarily and the number of possible broadcast schedules is infinite, it is not possible to find the broadcast schedule which gives the minimum average waiting time in all schedules when using search-based methods. Hence, in Section IV, for the case of dividing data into two segments of equal size, we present a schedule that gives the minimum average waiting time. In Section V we propose a heuristic scheduling method for dividing data into more than two segments, since many combinations of scheduling segments makes the analysis difficult.

#### IV. Dividing Data into Two Segments

Here we analyze the division of data into two segments of equal size. In this case, we find the schedule that gives the minimum average waiting time.
TABLE I
BROADCAST SCHEDULES WHICH INCLUDE 4 OR 5 SEGMENTS IN ONE
BROADCAST CYCLE (\(a = 5, D = 5\) MIN, \(N = 2\)).

<table>
<thead>
<tr>
<th>Broadcast Schedule</th>
<th>Average Waiting Time (sec)</th>
<th>Broadcast Schedule</th>
<th>Average Waiting Time (sec)</th>
</tr>
</thead>
<tbody>
<tr>
<td>(S_1S_2S_3)</td>
<td>23</td>
<td>(S_1S_2S_3S_4)</td>
<td>21</td>
</tr>
<tr>
<td>(S_1S_2S_2S_3)</td>
<td>38</td>
<td>(S_1S_2S_3S_2)</td>
<td>33</td>
</tr>
<tr>
<td>(S_1S_2S_3)</td>
<td>60</td>
<td>(S_1S_2S_3S_2)</td>
<td>51</td>
</tr>
<tr>
<td>(S_1S_2S_3)</td>
<td></td>
<td>(S_1S_2S_3S_2)</td>
<td>51</td>
</tr>
<tr>
<td>(S_1S_2S_3)</td>
<td></td>
<td>(S_1S_2S_3S_2)</td>
<td>33</td>
</tr>
<tr>
<td>(S_1S_2S_3)</td>
<td></td>
<td>(S_1S_2S_3S_2)</td>
<td>75</td>
</tr>
</tbody>
</table>

Let \(\alpha\) denote the integral part of \(a\), namely \(\alpha = [a]\). We define \(A\) by the following equation:

\[ A = (\alpha + 1)(\alpha - \alpha) - \alpha. \] (2)

The following statements give the broadcast schedule that allows the minimum average waiting time:

- For the case where \(A \leq 0\):

  \[ \frac{\alpha}{S_1 \cdots S_1 S_2}, \] (3)

  and the average waiting time is given by

  \[ \frac{\alpha + 3}{4a(a + 1)} \cdot D. \] (4)

- For the case where \(A \geq 0\):

  \[ \frac{\alpha + 1}{S_1 S_1 \cdots S_1 S_1 S_2}, \] (5)

  and the average waiting time is given by

  \[ \frac{3\alpha - 2\alpha + 6}{2a(a + 2)} \cdot D. \] (6)

V. DIVIDING DATA INTO MORE THAN TWO SEGMENTS

We propose a heuristic scheduling method of the division based broadcasting scheme, the segment insertion method, for dividing data into more than two segments.

A. Segment Insertion Method

The broadcast schedule \(A\) is produced according to the following process:

1. Let \(A\) be empty.
2. Divide the continuous media data into \(N\) segments \((S_1, \ldots, S_N)\) of equal size.
3. Add \(n_N\) pieces of \(S_N\) to \(A\).
4. First, let \(i = N - 1\) and add \(n_i\) pieces of \(S_i\) to the left side of each segment included in the schedule \(A\).
5. Decreasing \(i\), repeat step 2 until \(i = 1\).

With the segment insertion method, the number of pieces to be added, \(n_i\) \((i = 1, \cdots, N)\), is given by the following equations:

- For the case where \(i = 1\) and \(N = 2\):

  \[ n_1 = [a]. \] (7)

- For the case where \(i = 1\) and \(N = 3, 4, \cdots\):

  \[ n_1 = \left[ \frac{(N - 2)a + 1}{2N - 2} \right] - 1. \] (8)

- For the case where \(i = 2, \cdots, N\):

  \[ n_i = 1. \] (9)

Fig. 5 shows a process of the segment insertion method for \(a = 4\) and \(N = 4\).

B. Comparison with the Simple Repetition Scheme

Fig. 6 shows the average waiting time of the segment insertion method. Since the average waiting time is proportional to the playing time \(D\), the vertical axis is the average waiting time divided by \(D\), while the horizontal axis is the playback ratio \(a\). In addition, \(SI (N = i)' (i = 2, 3, 4)\) represents the average waiting time of the segment insertion method when the case that the data is divided into \(i\) segments, and \(N = i (c = i, \cdots, 10)'\) is the minimum average waiting time of all schedules including \(i\) to 10 segments in one broadcast cycle. Moreover, 'simple' means the average waiting time of the simple repetition scheme explained in Section II.

In the case of broadcasting an MP3-encoded 5-minute audio data with a 1-Mbps wireless broadcast systems such as satellite or bluetooth systems, \(a = 7.5\), as mentioned in Section III. Here, although the average waiting time of the simple
The server broadcasts the data, which is divided so that used. Let request data segment, and their waiting time is reduced. In intervals. Hence, clients have many chances to receive the first segment is shorter, the first segment is broadcast at frequent times are given arbitrarily.

This repetition scheme is 20 seconds, that of the 4-division segment insertion method is 8 seconds. This is 60% shorter.

VI. RELATED WORK

To reduce the waiting time under the continuity condition, various strategies have been studied so far.

In pyramid broadcasting (PB) [2], data, the quantity of which is $M$, is divided into $K$ segments, and $K$ channels are used. Let $D_i$ denote the playing time of the $i^{th}$ segment. The server broadcasts the $i^{th}$ segment of each data using the $i^{th}$ channel. Since the time required to broadcast precedent segments is shorter, the first segment is broadcast at frequent intervals. Hence, clients have many chances to receive the first segment, and their waiting time is reduced. In this scheme, however, the greater $K$ becomes, the bigger the data size of the $i^{th}$ segment is; therefore, buffer capacity requirements increase.

Skyscraper broadcasting (SB) [3] also considers this problem. Accordingly, a segment with a longer playing time than $D_i W$ is divided so that its playing time is $D_i W$. By modifying $W$, buffer requirements and the average waiting time are given arbitrarily.

Dynamic skyscraper broadcasting has also been proposed to deal with this problem [4]. This scheme is based on the SB scheme, and it dynamically changes broadcast data, and can improve the performance of SB when clients infrequently request data.

A permutation-based pyramid broadcasting [5] divides each channel into $p$ sub-channels, and a server broadcasts the $i^{th}$ segment using $p$ sub-channels. The beginning of the broadcast cycle of each sub-channel is shifted, and as a result, the waiting times are improved compared to those of PB. Although the paper suggests time division multiplexing which employs a single channel, the paper does not describe a detailed technique for reducing clients' waiting time via a single channel.

VII. CONCLUSION AND FUTURE WORK

In this paper, we proposed a division-based broadcasting scheme to reduce the waiting time of clients with a single channel. In our strategy, assuming that clients can play data continuously, the data is divided into several segments to produce an effective schedule. By broadcasting the first segment frequently, clients' waiting time is reduced. Thus, in this scheme, the number of segments and the broadcast schedule affect the waiting time.

For the case of dividing data into two segments, we also presented a schedule which gives the minimum average waiting time. The number of the first segment $S_1$ included in the broadcast schedule is determined by the ratio of the data playing time and the time needed to broadcast the segment, whereas the last segment $S_2$ is included only once. When the server broadcasts an MP3-encoded continuous media data using a 1-Mbps satellite system, the average waiting time is reduced by about 38% compared to the simple repetition scheme, which broadcasts data without division. Hence, dividing the data into only two segments is rather realistic.

For the case of dividing data into more than two segments, we proposed and evaluated the segment insertion method. As a result of analyzing the average waiting time of the segment insertion method, we found that the average waiting time is always shorter than or equal to that of the simple repetition scheme. The greater the number of segments $N$ is, the shorter the average waiting time becomes.

In future, we plan to study the delivery of multiple data with a single channel and other effective heuristics.

ACKNOWLEDGEMENT

This research was supported in part by "The 21st Century Center of Excellence Program" of the Ministry of Education, Culture, Sports, Science and Technology, Japan, and Special Coordination Funds for Promoting Science and Technology of the Ministry of Education, Culture, Sports, Science and Technology of Japan.

REFERENCES