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Author(s)	Hashizume, Yoko
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Hashizume, Yoko  
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## ***On the Uniqueness of the Decomposition of a Link***

By Yoko HASHIZUME

### **§ 0. Introduction.**

A link of multiplicity  $n$  is a collection of  $n$  disjoint simple closed oriented polygons in the 3-sphere  $S^3$ . Especially a link of multiplicity 1 is a so-called knot. H. Schubert [1] showed that the genus of the product of two knots is equal to the sum of their genera and that every knot is decomposable in a unique way into prime knots. The purpose of this paper is to extend his results to the case of links.

In § 1, using some of the results and methods due to H. Schubert [1], we define the product of links and prove some preliminary theorems.

In § 2, we define decomposition systems for non trivial and non separable links, and prove by the aid of the decomposition system the following

**MAIN THEOREM.** *Every non trivial and non separable link can be decomposed uniquely into prime links.*

For the links of multiplicity 1, this theorem coincides with H. Schubert's result. But our proof is simpler than his.

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**§ 1.** A link of multiplicity  $n$  is a collection of  $n$  disjoint simple closed oriented polygons in the 3-sphere  $S^3$ <sup>1)</sup>. Two links  $l$  and  $l'$  are said to be *equivalent* and denoted by  $l \approx l'$ , if there exists an orientation preserving semilinear mapping  $S^3$  onto itself which maps one of them onto the other. Especially, a link of multiplicity 1 is a so-called knot. Throughout this paper we shall denote by  $l$  a link, and by  $k$  a knot.

We shall say that  $l$  has  $\mu$  components, if there are  $\mu$  disjoint cubes  $Q_1, \dots, Q_\mu$  for  $l$  such that  $l \cap Q_i = \emptyset$ ,  $l \cap Q_i \neq \emptyset$  ( $i = 1, \dots, \mu$ ) and there are no  $\mu+1$  disjoint cubes with these properties. A link with multiplicity

1) In the following  $S^3$  will be taken as the boundary of a 4-simplex in 4-dimensional Euclidean space  $R^4$ . To simplify our observation, we chose "infinity" of  $S^3$  as a vertex of this 4-simplex, and call the opposite 3-simplex the base simplex and denote it by  $A^3$ .

2) Such expressions as cubes, spheres, surfaces, disks, arcs, etc. should be understood all simplicial and mappings should be understood all semilinear  $S^3$  onto itself.

1 is said to be *non separable*. A link  $l$  is said to be *trivial*, if each component of  $l$  consists of a single triangle.

Let  $Q$  be a solid cube and let  $l=(k_1, \dots, k_n)$  be a link which has an arc  $v$  of  $k_i$  in common with the boundary  $\dot{Q}$  of  $Q$ , the remaining  $l-v$  lying wholly within  $Q$  except for  $v$ . Setting  $\overline{l-v}=l^*$ , we shall say that  $(l^*, Q)$  is a *representation of  $l$  (with respect to  $k_i$ )*, or  $l^*$  represents  $l$  in  $Q$  (*with respect to  $k_i$* ). We shall call  $v$  a *joining arc of  $l^*$  on  $\dot{Q}$* . If we replace  $v$  by an arbitrary joining arc  $v'$  of  $l^*$  on  $\dot{Q}$  such that  $l^* \cup v'$  makes a link  $l'$ , then we have  $l' \approx l$  by an analogous argument to Satz 1 of [1].

A cube  $Q$  will be said to be an *admissible cube* of  $l$ , if  $l \cap \dot{Q}$  consists of just two points. We shall denote by  $Q^c$  the complementary cube of  $Q$ , i.e.  $\overline{S^3 - Q}$ .

Let  $Q$  be an admissible cube of  $l$ . If  $l \cap Q$  represents  $l_1=(k_1^1, \dots, k_n^1)$  in  $Q$  with respect to  $k_i^1$  and  $l \cap Q^c$  represents  $l_2=(k_1^2, \dots, k_m^2)$  in  $Q^c$  with respect to  $k_j^2$ , then  $l$  is said to be a *product of  $l_1$  and  $l_2$  associated with  $(k_i^1, k_j^2)$* , and denoted by  $l=l_{i,j} \cdot l_2$ . If we take no notice of the locality of the product, then we say merely that  $l$  is a *product of  $l_1$  and  $l_2$*  and denote  $l$  by  $l_1 \cdot l_2$ . We say also that  $Q$  *cut out  $l_1$  from  $l$*  and  $l_2$  is the *rest link of  $l$  with respect to  $Q$* . Then we have

**Lemma 1.** *For every two links  $l=(k_1^1, \dots, k_n^1)$  and  $l_2=(k_1^2, \dots, k_m^2)$ , there exists  $l_{i,j} \cdot l_2$  uniquely, where  $i$  and  $j$  are a given pair of integers with  $1 \leq i \leq n$  and  $1 \leq j \leq m$ .*

The proof is essentially the same as Satz 3 of [1] and is omitted.

We shall denote by  $m(l)$  the multiplicity of  $l$ . Then it follows directly from the definition of the products:

**Theorem 1.** *If  $l=l_1 \cdot l_2$ , then  $m(l)=m(l_1)+m(l_2)-1$ .*

$l$  is said to be *prime*, if  $l$  is non separable and if, whenever  $l=l_1 \cdot l_2$ , one of  $l_1$ ,  $l_2$  is trivial.

#### EXAMPLES.

(1)  $l_1$  (Fig. 1) is a prime link.

(2)  $l_2$  (Fig. 2) is a non prime link which is the product of two  $l_1$ .

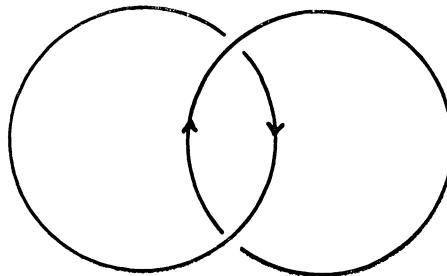
In the same way as the case for knots [2], we can span for an arbitrary given link  $l$  a connected orientable singularity free surface  $F$ . If we denote by  $g(F)$  the genus of  $F$ , the minimal number of  $g(F)$  for all choices of such  $F$  will be called the *genus of  $l$*  and will be denoted by  $g(l)$ . Then we have the following

**Theorem 2.** *If  $l=l_1 \cdot l_2$ , then  $g(l)=g(l_1)+g(l_2)$ .*

REMARK. A product  $l = l_{i,j} \cdot l_2$  depends in general on the pair  $(i, j)$ , but its genus is uniquely determined only by  $l_1$  and  $l_2$ .

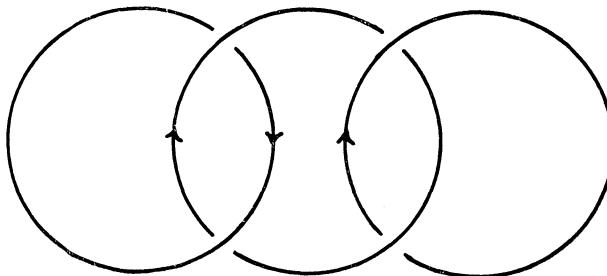
Since Theorem 2 is proved by an analogous way to the proof of Satz 4 of [1], using a lemma which is a generalization of Hilfssatz 7 of [1] for the case of links, the proof is omitted.

For the case of knots, we know that  $g(k) = 0$ , if and only if  $k$  is trivial and that if  $g(k) = 1$  then  $k$  is prime (cf. Satz 6 of [1]). But by



$l_1$

Fig. 1.



$l_2$

Fig. 2.

our definition of links, the circumstances are very different from the case of knots. For example,  $l_1$  (Fig. 1) is a link of genus 0 but non trivial, and  $l_2$  (Fig. 2) is a link of genus 0 but neither trivial nor prime. We can only verify the following :

**Theorem 3.** *If  $g(l) = 0$  and  $m(l) = 2$ , and if  $l$  is non separable, then  $l$  is prime.*

Proof. Let  $l = l_1 \cdot l_2$ . Then by Theorem 2  $g(l_1) = g(l_2) = 0$ . On the

other hand, since  $m(l)=2$ , either  $m(l_1)=1$  or  $m(l_2)=1$ . We may assume without loss of generality that  $m(l_1)=1$ . Then from  $g(l_1)=0$  it follows that  $l_1$  is trivial. Hence  $l$  is prime.

We have further

**Theorem 4<sup>3)</sup>.** *Let  $l=l_1 \cdot l_2$ . If one of  $l_1, l_2$  is a trivial knot, then the other is equivalent to  $l$ . Conversely if one of  $l_1, l_2$  is equivalent to  $l$ , then the other is a trivial knot.*

Proof. The first part is clear (cf. Satz 5 of [1]). To prove the second part, let  $l \approx l_1$ . Then clearly  $m(l)=m(l_1)$ . But since  $m(l)=m(l_1 \cdot l_2)=m(l_2)+m(l_2)-1$ , we have  $m(l_2)=1$ , what is the same,  $l_2$  is a knot. From Theorem 2, it follows that  $g(l_2)=0$ , i.e.  $l_2$  is a trivial knot.

**§ 2.** In this section we shall consider only about non separable links.

We shall use the following notations. Let  $Q$  be the admissible cube of  $l$ , that is, a cube intersecting  $l$  in just two points, and let  $Q_1, \dots, Q_n$  be admissible cubes of  $l$  in  $\text{Int } Q$  such that  $\dot{Q}_i \cap \dot{Q}_j = \emptyset$  ( $i, j = 1, \dots, n$ ). Then take out all the maximal cubes<sup>4)</sup>  $\{Q_{n1}, \dots, Q_{nm}\}$  from  $\{Q_1, \dots, Q_n\}$ . Let  $v, v_{n1}, \dots, v_{nm}$  be joining arcs of  $l$  on  $\dot{Q}, \dot{Q}_{n1}, \dots, \dot{Q}_{nm}$  respectively. Then  $l \cap (S^3 - (Q^c \cup Q_{n1} \cup \dots \cup Q_{nm}))$  together with  $v, v_{n1}, \dots, v_{nm}$  make a link as shown in Fig. 3. We denote it by  $l - (Q^c + Q_1 + \dots + Q_n)$ .

**Lemma 2.** *Let  $l=l_1 \cdot l_2$  and let  $(l^*, Q)$  be a representation of  $l$ . Then there exists an admissible cube  $Q_1$  of  $l$  inside  $Q$  such that  $l_1^* = l^* \cap Q_1$  represents  $l_1$  in  $Q_1$  with respect to  $k_i$  and  $l - (Q^c + Q_1)$  is equivalent to  $l_2$  (Fig. 4).*

Since Lemma 2 is easily proved almost in the same way as the proof of Hilfssatz 8 of [1], we omit the proof.

$\mathfrak{D} = \{(l_1, Q_1), \dots, (l_n, Q_n)\}$  will be called a *decomposition system of  $l$* , if it satisfies the following conditions :

- (D I) Each  $Q_i$  is an admissible cube of  $l$ .
- (D II) For each pair of different  $Q_i$  and  $Q_j$ ,  $Q_i \cap Q_j = \emptyset$  or  $Q_i \subset \text{Int } Q_j$  or  $Q_j \subset \text{Int } Q_i$ , or  $Q_i = Q_j$ .
- (D III)  $S^3 = Q_1 \cup \dots \cup Q_n$ .
- (D IV) Let  $Q_j, \dots, Q_i$  be all the cubes contained in  $\text{Int } Q_i$ . Then  $l_i \approx l - (Q^c + Q_j + \dots + Q_i)$  and each  $l_i$  is prime.
- (D V)  $l$  is a product of  $l_1, \dots, l_n$ .

3) Theorem 4 is the extension of Satz 5 of [1] to links.

4) We say that  $Q_i$  is maximal in  $\{Q_1, \dots, Q_n\}$ , if there is no  $Q_j$  in  $\{Q_1, \dots, Q_n\}$  such that  $Q_i \subset \text{Int } Q_j$ .

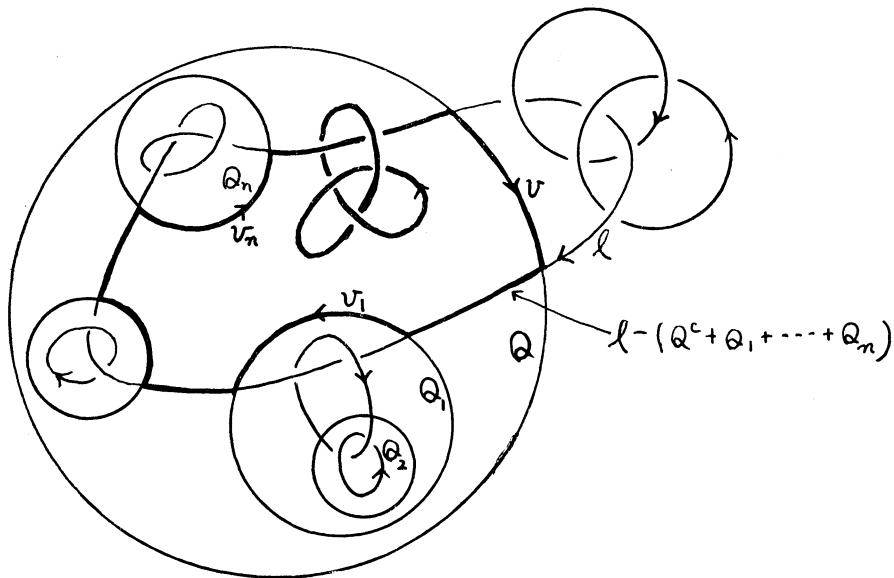


Fig. 3.

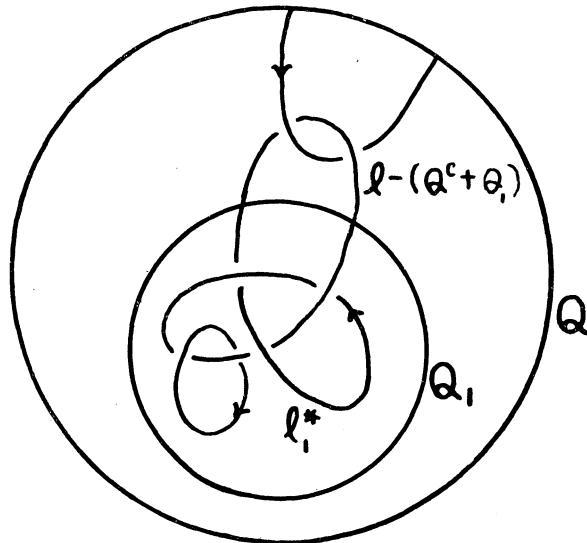


Fig. 4.

Now we shall show that

**Theorem 5.** *For every non trivial and non prime link, there exists at least one decomposition system.*

Proof. Since  $l$  is neither trivial nor prime, there exists an admissible cube  $Q_1$  such that  $l = l_1 \cdot l_2$ , where  $l_1^* = l \cap Q_1$ ,  $l_2^* = l \cap Q_1^c$ .

From Theorem 1 and Theorem 2 we have

$$(m(l_1) + g(l_1)) + (m(l_2) + g(l_2)) = m(l) + g(l) + 1$$

Here  $m(l_i) \geq 1$  ( $i = 1, 2$ ). But if e.g.  $m(l_1) = 1$  and  $g(l_1) = 0$ , then  $l_1$  must be trivial, contradicting the way of decomposition, therefore  $m(l_1) + g(l_1) \geq 2$ . Thus we have

$$m(l_i) + g(l_i) < m(l) + g(l) \quad i = 1, 2 \quad (1).$$

Here we put  $Q_2 = Q^c$ . If  $l_i$  ( $i = 1, 2$ ) is not prime, then there exists by Lemma 2 an admissible cube  $Q_{ii}$  in  $\text{Int } Q_i$  such that  $l_i = l_{i1} \cdot l_{i2}$ , where  $l_{i1}^* = l_i \cap Q_{ii}$ ,  $l_{i2} = l_i - (Q_i^c + Q_{ii})$ . Next suppose  $l_{i2}$  ( $i = 1, 2$ ) is not prime.

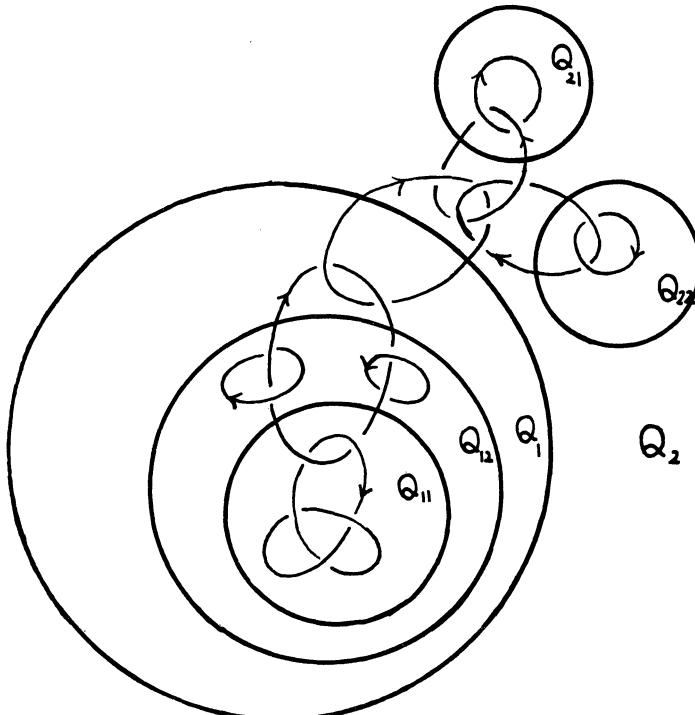


Fig. 5.

Then we decompose  $l_{i2}$  into two links with the aid of an admissible cube  $Q_{i2}$  of  $l$  lying wholly within  $Q_i$  and such that  $Q_{i1} \cap Q_{i2} = \emptyset$  as shown in Fig. 5.

We proceed with these decomposition procedures. Then these procedures cannot continue indefinitely as will be seen easily from (1), and each factor link comes finally to be prime.

Renumbering the cubes and links above obtained, we obtain a system  $\mathfrak{D} = \{(l_1, Q_1), \dots, (l_n, Q_n)\}$ . From the way of the construction of  $\mathfrak{D}$  it is clear that  $\mathfrak{D}$  satisfies the conditions (D I)–(D V). Thus Theorem 5 is proved.

Let  $\mathfrak{D} = \{(l_1, Q_1), \dots, (l_n, Q_n)\}$  and  $\mathfrak{D}' = \{(l'_1, Q'_1), \dots, (l'_m, Q'_m)\}$  be decomposition systems of a link. An element  $(l_i, Q_i)$  of  $\mathfrak{D}$  and an element  $(l'_j, Q'_j)$  of  $\mathfrak{D}'$  will be called *equivalent* and denoted by  $(l_i, Q_i) \approx (l'_j, Q'_j)$ , if  $l_i \approx l'_j$ .  $\mathfrak{D}$  and  $\mathfrak{D}'$  will be called *equivalent* and denoted by  $\mathfrak{D} \approx \mathfrak{D}'$ , if  $n=m$  and if there is a one-to-one correspondence between the elements of  $\mathfrak{D}$  and the elements of  $\mathfrak{D}'$  such that the corresponding elements are equivalent.  $(l_i, Q_i)$  will be called *an end of  $\mathfrak{D}$* , if there is no  $Q_j$  in  $\mathfrak{D}$  such that  $Q_j \subset \text{Int } Q_i$ .

**Lemma 3.** *For arbitrary given two decomposition systems  $\mathfrak{D} = \{(l_1, Q_1), \dots, (l_n, Q_n)\}$  and  $\mathfrak{D}' = \{(l'_1, Q'_1), \dots, (l'_m, Q'_m)\}$  of the same  $l$  and an end of  $\mathfrak{D}'$ , say  $(l'_1, Q'_1)$ , there exists a decomposition system  $\mathfrak{D}'' = \{(l''_1, Q''_1), \dots, (l''_n, Q''_n)\}$  of  $l$  satisfying the following conditions:*

- (1)  $\mathfrak{D}'' \approx \mathfrak{D}$ .
- (2)  $(l''_1, Q''_1)$  is an end of  $\mathfrak{D}''$  such that  $l''_1 \approx l'_1$  and either  $l''_1 \subseteq l'_1$  or  $l''_1 \supseteq l'_1$ .

**Proof.** First we deform  $l$ ,  $\mathfrak{D}$  and  $\mathfrak{D}'$  by an isotopic simplicial deformation into the following forms without changing the other circumstances<sup>5)</sup>:

- (1)  $Q'_1$  is a 3-simplex in the interior of the base simplex  $\Delta^3$ .
- (2) Let  $l \cap Q'_1 = \{A, B\}$ . Then  $A$  and  $B$  are respectively in the interior of certain 2-simplexes which are boundary simplexes of  $Q'_1$ .
- (3)  $l, \{\dot{Q}_1, \dots, \dot{Q}_n\}, \{\dot{Q}'_1, \dots, \dot{Q}'_m\} \subset \text{Int } \Delta^3$ .
- (4) Neither  $A$  nor  $B$  is a vertex of  $l$ .

Thereafter we give a simplicial decomposition  $\mathfrak{B}$  on  $l \cup (\bigcup_i \dot{Q}_i)$  such that  $\mathfrak{B}$  is so fine that if  $\sigma$  is a simplex of  $\mathfrak{B}$  such that  $\sigma \cap Q'_1 \neq \emptyset$ , then there is no simplex  $\sigma'$  of  $\mathfrak{B}$  such that  $\sigma' \cap \{(\bigcup_i \dot{Q}_i) - Q'_1\} \neq \emptyset$  and  $\sigma' \cap \sigma \neq \emptyset$ .

From the above condition and (1)–(4), we can further deform  $Q'_1$  and  $\mathfrak{D}$  such that it satisfies the following conditions, without moving  $l$  and without changing the other circumstances.

5) See pp. 89–pp. 90 [1].

- (5)  $A, B \notin \dot{Q}_i$  ( $i = 1, \dots, n$ ).
- (6) Neither  $A$  nor  $B$  is a vertex of  $\mathfrak{B}$ .
- (7) Every vertex of  $Q'_i$  is not on  $\dot{Q}_i$  ( $i = 1, \dots, n$ ).
- (8) The intersection of any edge of  $Q'_i$  and any edge of  $\mathfrak{B}$  is null.

From the above conditions it follows that  $\dot{Q}'_i \cap (\cup \dot{Q}_i)$  consists of a finite number of simple closed curves. We call them the *cut lines on  $\dot{Q}'_i$* . By the usual method, we are going to delete these cut lines from  $\dot{Q}'_i$ , replacing  $\mathfrak{D}$  by an equivalent one. Take out one of the innermost cut lines on  $\dot{Q}'_i$ , and name it  $c$ . Then there is a disk  $E$  on  $\dot{Q}'_i$  bounded by  $c$  such that there are no more cut lines on  $E$ . First we are going to delete  $c$  without getting new cut lines on  $\dot{Q}'_i$ . Suppose  $c \subset \dot{Q}_i \cap \dot{Q}'_i$ . Since  $\dot{Q}_i \cap E = c$  the following two cases are to be considered:

Case I:  $E \subset Q_i$  and Case II:  $E \subset Q_i^c$ .

Case I.  $c$  divides  $\dot{Q}_i$  into two disks  $D_1, D_2$  such that  $D_1 \cap D_2 = \dot{D}_i = \dot{D}_2 = c$  and  $D_1 \cup D_2 = \dot{Q}_i$ .  $E$  divides  $Q_i$  into two cubes  $Q_{i1}, Q_{i2}$  such that  $E = Q_{i1} \cap Q_{i2}$ ,  $Q_{i1} \cup Q_{i2} = Q_i$  and  $\dot{Q}_{ij} = D_j \cup E$  ( $j = 1, 2$ ). We put hereafter  $l \cap \dot{Q}_i = \{C, D\}$ .

The following three cases actually occur in this case.

- (0) Neither  $A$  nor  $B$  is contained in  $E$ .
- (1) One of the  $A, B$  is contained in  $E$  and the other is not contained in  $E$ .
- (2) Both  $A$  and  $B$  are contained in  $E$ .

Case I (0).  $E \subset Q_i$ .  $A, B \notin E$ .

In this case, either  $C, D \in D_1$  or  $C, D \in D_2$ , because  $l \cap \dot{Q}_{ij} = l \cap D_j$  ( $j = 1, 2$ ) must consist of even number of points. Without loss of generality we may assume that  $C, D \in D_1$  as Fig. 6.

Then  $l \cap Q_{i2} = \emptyset$ . For, if  $l \cap Q_{i2} \neq \emptyset$ , then from the facts  $l \cap \dot{Q}_{i2} = \emptyset$  and  $l \cap Q_{i2}^c \neq \emptyset$  ( $\exists C, D$ ), it follows that  $l$  is separable, which contradicts our first assumption.

Hence  $l \cap Q_{i2} = \emptyset$ . Therefore  $l \cap Q_i = l \cap Q_{i1}$ . Then we cut off  $Q_{i2}$  from  $Q_i$ . Shifting slightly all the vertices of a certain simplicial decomposition of  $\dot{Q}_i$  on  $E$  to one direction, we can delete the cut line  $c$  from  $\dot{Q}'_i$  as shown in Fig. 7, without giving rise to any new cut lines on  $\dot{Q}'_i$  and without changing the other circumstances. We write the cube thus obtained still by the same symbol  $Q_{i1}$ . Then replacing  $(l_i, Q_i)$  by  $(l_i, Q_{i1})$  in the decomposition system  $\mathfrak{D}$ , we obtain a decomposition system  $\tilde{\mathfrak{D}}$  which is obviously equivalent to  $\mathfrak{D}$ . Since  $c$  is deleted, the number of

6) Such simplicial decomposition  $\mathfrak{B}$  is obtained by refining arbitrary given one.

the cut lines on  $Q'_i$  is diminished at least by one from the original one. If  $C, D \in D_2$ , then we may merely exchange the number 1 and 2 in the above process.

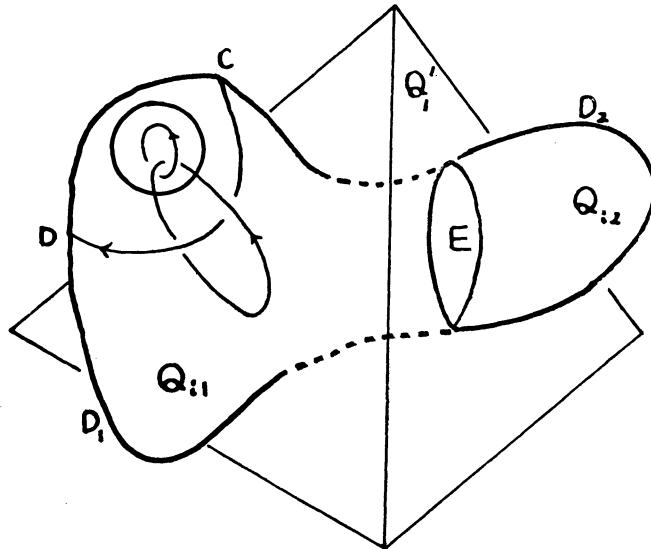


Fig. 6.

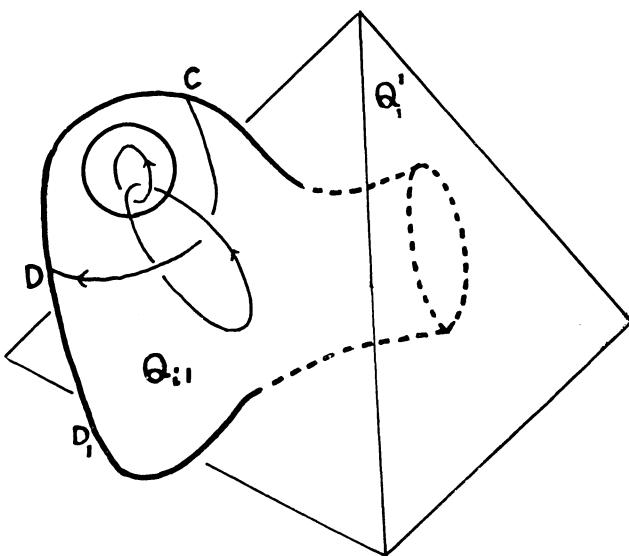


Fig. 7.

Case I (1).  $E \subset Q_i$ .  $A \in E, B \notin E$ .

In this case one of  $C, D$  is contained in  $D_1$  and the other is contained

in  $D_2$ . We may assume without loss of generality that  $C \in D_1$  and  $D \in D_2$ . By the definition of the decomposition system  $l_i = l - (Q^c + Q_j + \dots + Q_l)$ , where  $Q_j, \dots, Q_l$  are all the cubes of  $\mathfrak{D}$  in  $\text{Int } Q_i$ . Let  $l_i^* = \overline{l_i - v_i}$ , where  $v_i$  is the joining arc of  $l$  on  $Q_i$  and let  $l_i^* \cap Q_{i1}, l_i^* \cap Q_{i2}$  represent  $l_{i1}, l_{i2}$  in  $Q_{i1}, Q_{i2}$  respectively. Then it is easy to see that  $l_i = l_{i1} \cdot l_{i2}$ . Since  $l_i$  is prime, one of  $l_{i1}, l_{i2}$  must be trivial. We may assume without loss of generality that  $l_{i2}$  is trivial. Then by Théorem 4,  $l_i \approx l_{i1}$  (Fig. 8).

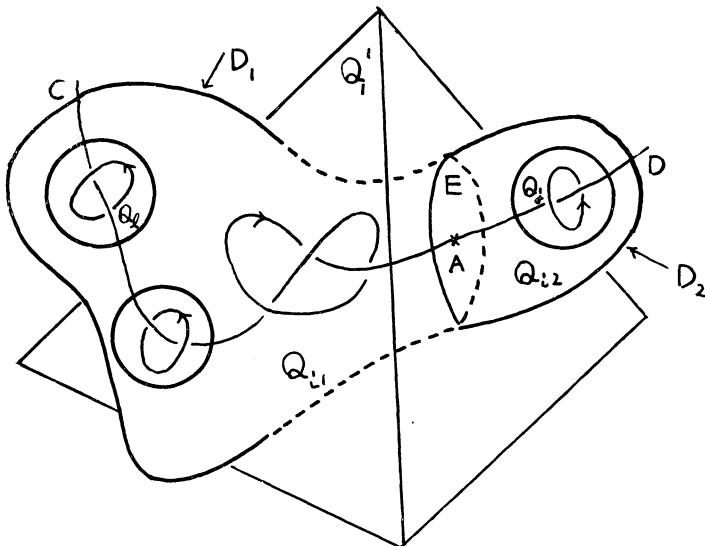


Fig. 8.

Then we cut off  $Q_{i2}$  from  $Q_i$ . Thereafter deforming  $Q_{i1}$  in the same way as Case I (0), we delete the cut line  $c$  (Fig. 9). Replacing  $(l_i, Q_i)$  by  $(l_{i1}, Q_{i1})$  in  $\mathfrak{D}$ , we obtain a decomposition system  $\tilde{\mathfrak{D}}$  which is obviously equivalent to  $\mathfrak{D}$ .

Case I (2).  $E \subset Q_i$ .  $A, B \in E$ .

In this case, it follows without loss of generality that  $C, D \in D_1$ . Then we may assume that the joining arc  $v_i$  of  $l_i^*$  on  $Q_i$  is contained in  $D_1$ . Let  $l \cap Q_{i2}$  represent  $l_{i2}$  in  $Q_{i2}$  and let the rest link of  $l_i$  with respect to  $Q_{i2}$  be  $l_{i1}$ . Then  $l_i = l_{i1} \cdot l_{i2}$  as shown in Fig. 10. Since  $l_i$  is prime, one of  $l_{i1}, l_{i2}$  must be trivial.

First suppose that  $l_{i_1}$  is trivial. Then we cut off  $Q_{i_1}$  from  $Q_i$ . Deforming then  $Q_{i_2}$  by the usual way, we may delete the cut line  $c$ . Since  $l_{i_1}$  is trivial,  $l_{i_2} \approx l_i$ . Moreover  $l \cap Q_{i_2}$  clearly consists of just two points  $A, B$ . Replacing then  $(l_i, Q_i)$  by  $(l_{i_2}, Q_{i_2})$  in  $\mathfrak{D}$ , we obtain a decomposition system  $\tilde{\mathfrak{D}}$  which is obviously equivalent to  $\mathfrak{D}$ .

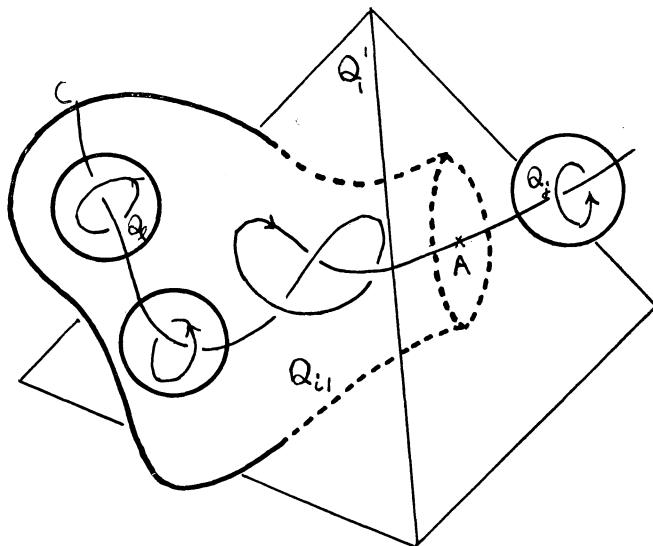


Fig. 9.

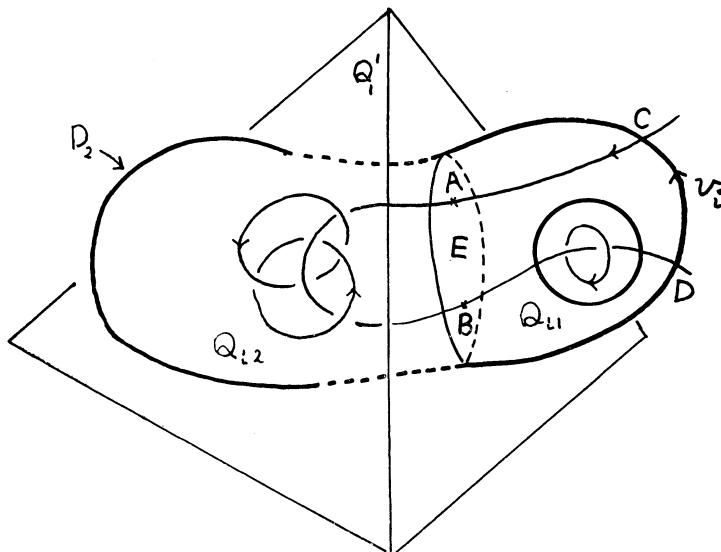


Fig. 10.

Next suppose that  $l_{i2}$  is trivial. It is easy to see that either  $l \cap Q_{i2} \subset Q'_1$  or  $l \cap Q_{i2} \subset Q'^c_1$ .

Suppose that  $l \cap Q_{i2} \subset Q'^c_1$ . Then  $l - (l \cap Q_{i2}) \subset Q'_1$ . Let  $u$  be an arc joining  $A$  and  $B$  on  $E$  as shown in Fig. 11. Then since  $(l'_1, Q'_1)$  is an end of  $\mathcal{D}'$ ,  $(l - (l \cap Q_{i2})) \cup u \approx l'_1$ . Set  $(l - (l \cap Q_{i2})) \cup u = l'$  and  $l' \cap Q_i = l_{i1}$

and let the rest link of  $l'$  with respect to  $Q_i$  be  $l''$ . Then obviously  $l' = l_{i_1} \cdot l''$ . Since  $l'_i$  is prime and since  $l_{i_1}$  is not trivial by assumption,  $l''$  must be trivial. Then  $\mathfrak{D}$  does not satisfy the condition (D IV), which contradicts that  $\mathfrak{D}$  is a decomposition system of  $l$ .

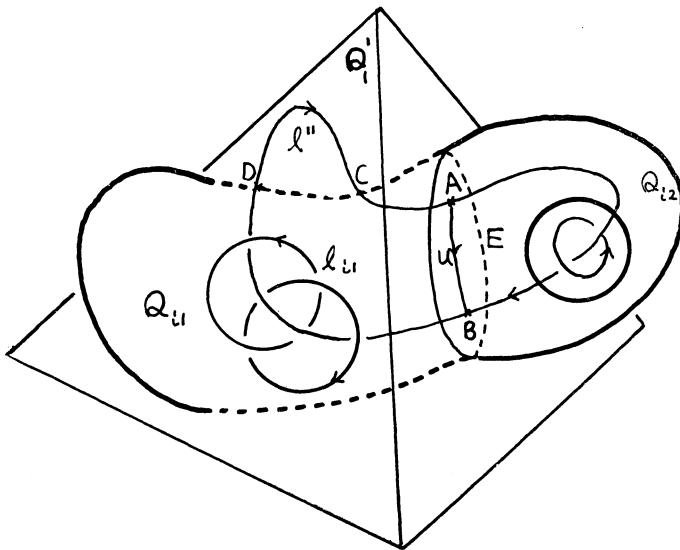


Fig. 11.

Hence we have  $l \cap Q_{i2} \subset Q'_i$ . Since  $(l'_i, Q'_i)$  is an end of  $\mathfrak{D}'$  and  $l_{i2}$  is trivial, one and only one cube of  $\mathfrak{D}$ , say  $Q_i$ , is contained in  $Q_{i2}$  as shown in Fig. 12. Then it is easy to see that  $(l_i, Q_i)$  is an end of  $\mathfrak{D}$  and  $l_i \approx l'_i$  and  $l_i^* \subset l'_i^*$ . When  $l_{i2}$  is trivial in Case I (2), thus Lemma 3 is proved.

Case II.  $E \subset Q_i^c$ .

$E$  divides  $Q_i^c$  into two cubes  $M_1, M_2$  such that  $M_1 \cap M_2 = E$ ,  $M_1 \cup M_2 = Q_i^c$  and  $M_j = E \cup D_j$  ( $j = 1, 2$ ). Since there is no cut line on  $E$ , one of  $M_1$  or  $M_2$ , say  $M_1$ , is contained in  $Q'_i$ .

Let  $E, D_1, D_2$ , etc. have the same meanings as in Case I. There are three cases, Case II (0), Case II (1) and Case II (2) to be considered.

Case II (0).  $E \subset Q_i^c$ .  $A, B \notin E$ .

If  $l \cap M_1 = \emptyset$ , then replace  $Q_i$  by  $\tilde{Q}_i = Q_i \cup M_1$ . Then we delete  $c$  by the usual way as shown in Fig. 13, and replace  $(l_i, Q_i)$  by  $(l_i, \tilde{Q}_i)$  in  $\mathfrak{D}$ .  $\tilde{\mathfrak{D}} = \{(l_1, Q_1), \dots, (l_i, \tilde{Q}_i), \dots, (l_n, Q_n)\}$  thus obtained is obviously a decomposition system of  $l$  equivalent to  $\mathfrak{D}$ . Therefore we assume below that  $l \cap M_1 \neq \emptyset$ .

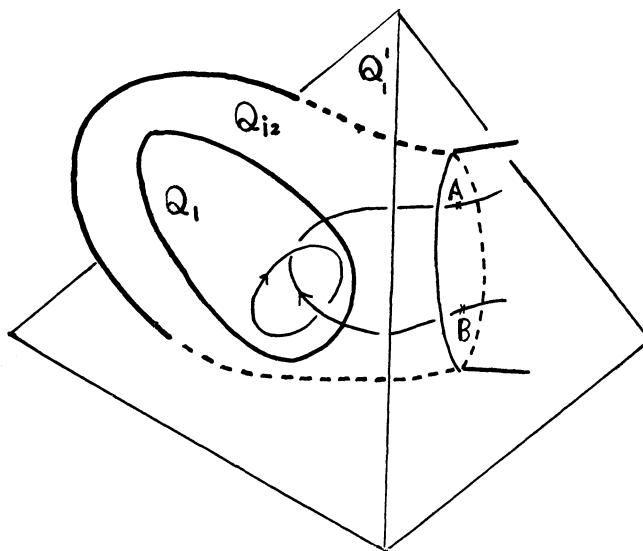


Fig. 12.

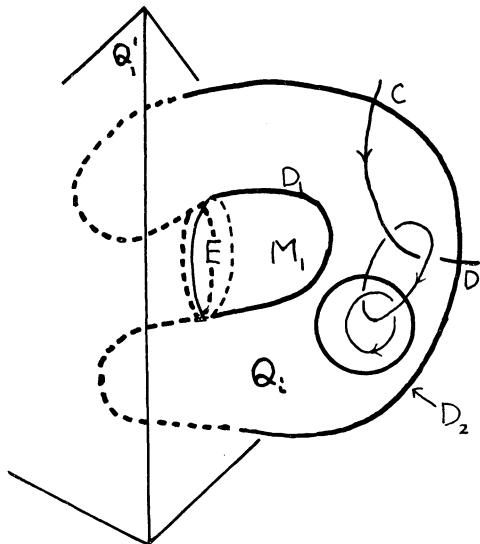


Fig. 13.

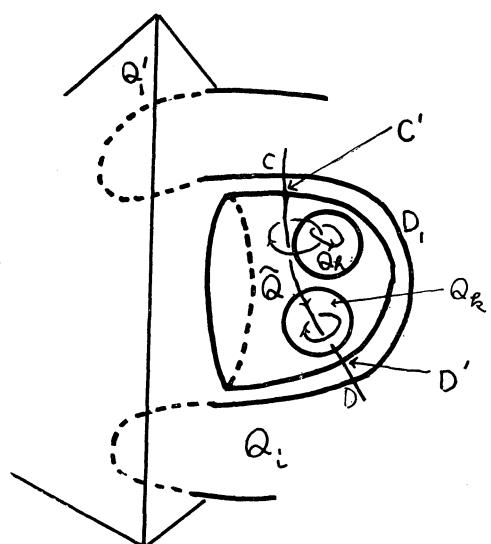


Fig. 14.

From the above assumption and from the fact that  $A, B \notin E$ , both  $C$  and  $D$  must be on  $D_1$  since  $l$  is non separable. We take a cube  $\tilde{Q}$  in  $\text{Int}M_1$  so close by  $M_1$  that  $l \cap \tilde{Q} = \{C', D'\}$  and  $CC', DD'$  are line segments and  $\tilde{Q} \cap Q_j = \emptyset$  ( $j = 1, \dots, n$ ) as shown in Fig. 14. Let  $Q_1, \dots, Q_k$  be all the cubes of  $\mathfrak{D}$  contained in  $\tilde{Q}$  (there may be no such cubes). It is

easy to see that  $l \subset M_i \cup Q_i$ . Since there is no cube in  $\mathfrak{D}$  which contains  $\tilde{Q}_i$ , one of  $Q_i^c, Q_h^c, \dots, Q_k^c$ , say  $Q_i^c$  must coincide with some cube of  $\mathfrak{D}$ , say  $Q_1$ . If we set  $l - (\tilde{Q}^c + Q_h + \dots + Q_k) = \tilde{l}$ , then clearly  $\tilde{l} \approx l_i$ . Further it is clear that  $l - (\tilde{Q} + Q_j + \dots + Q_l) \approx l_i$ , where  $Q_j, \dots, Q_l$  are all the cubes of  $\mathfrak{D}$  in  $\text{Int } Q_i$ . Then replacing  $(l_i, Q_i)$  by equivalent  $(\tilde{l}, \tilde{Q})$  and  $(l_i, Q_i)$  by equivalent  $(l - (\tilde{Q} + Q_j + \dots + Q_l), \tilde{Q}^c)$  in  $\mathfrak{D}$ , we obtain a decomposition system  $\tilde{\mathfrak{D}}$  of  $l$  which is equivalent to  $\mathfrak{D}$ . Thus the cut line  $c$  is deleted from  $\dot{Q}_i'$ , and from the way of the construction of  $\tilde{\mathfrak{D}}$ , there arise no new cut lines on  $\dot{Q}_i'$ .

Case II (1).  $E \subset Q_i^c$ .  $A \in E, B \notin E$ .

In this case, one of  $C, D$  is contained in  $D_1$  and the other is contained in  $D_2$ . We may assume without loss of generality that  $C \in D_1$ . We take a cube  $\tilde{Q}$  in  $\text{Int } M_1$  so close by  $M_1$  that  $l \cap \tilde{Q} = \{A', C'\}$  and  $AA', CC'$  and line segments and  $\tilde{Q} \cap \tilde{Q}_j = \emptyset$  ( $j=1, \dots, n$ ). Let  $Q_h, \dots, Q_k, Q_j, \dots, Q_l$  and  $\tilde{l}$  have the same meanings as in Case II (0).

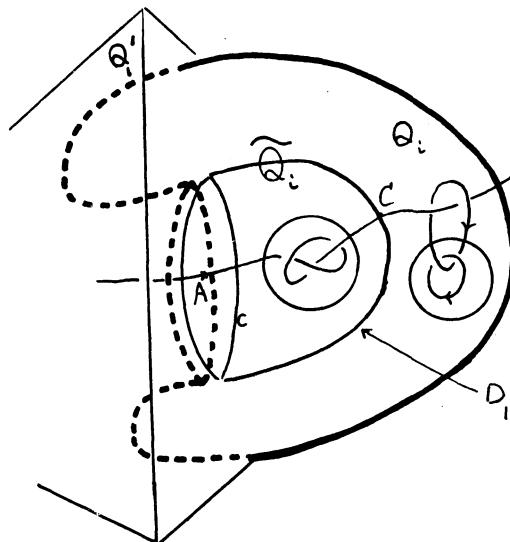


Fig. 15.

If  $\tilde{l}$  is trivial, then we replace  $Q_i$  by  $\tilde{Q}_i = Q_i \cup M_1$ . It is clear that  $\tilde{Q}_i$  is also an admissible cube of  $l$ . Then obviously  $l_i = l - (Q_i^c + Q_j + \dots + Q_l) \approx l - (\tilde{Q}_i^c + Q_j + \dots + Q_l + Q_h + \dots + Q_k)$ . We then delete  $c$  by the usual way as shown in Fig. 15. Then replacing  $(l_i, Q_i)$  by  $(l - (\tilde{Q}_i^c + Q_j + \dots + Q_l + Q_h + \dots + Q_k), \tilde{Q}_i)$  in  $\mathfrak{D}$ , we obtain a decomposition system  $\tilde{\mathfrak{D}}$  of  $l$  which is equivalent to  $\mathfrak{D}$ .

If  $\tilde{l}$  is non trivial, then the following two cases actually occur: (A) There exists a  $Q_r$  in  $\mathfrak{D}$  such that  $\tilde{Q}_i \subset \text{Int } Q_r$ . (B) There is no such  $Q_r$  as (A).

Case (A). In this case we choose  $Q_r$  such that there is no  $Q_t$  in  $\mathfrak{D}$  such as  $\tilde{Q}_i \subset Q_t \subset Q_r$ , as shown in Fig. 16. Let  $Q_i, Q_j, \dots, Q_s$  be all the

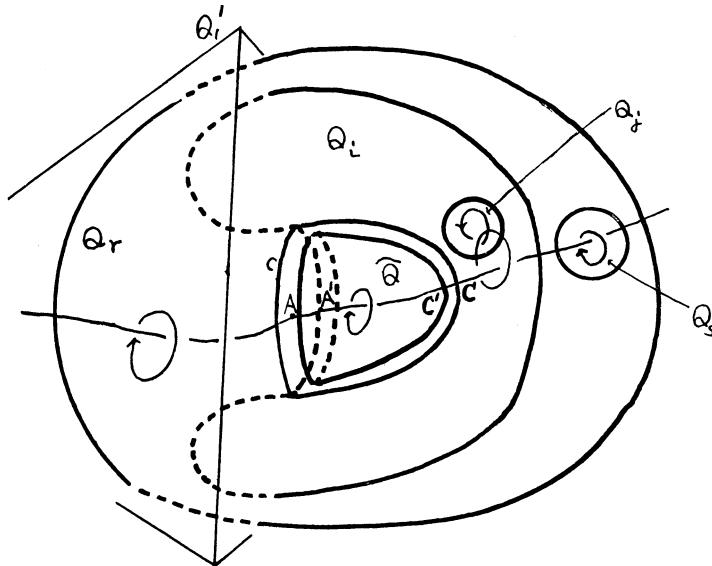


Fig. 16.

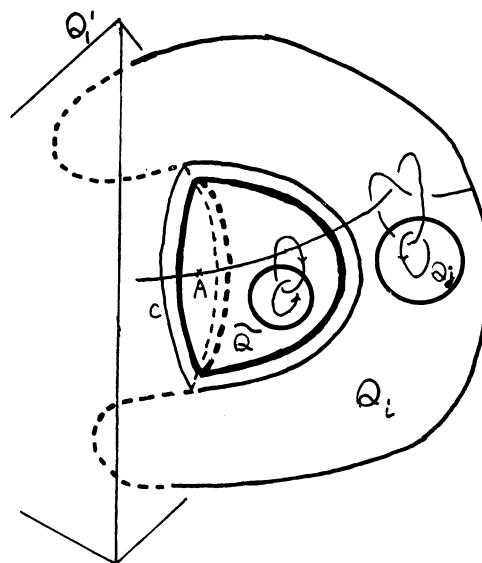


Fig. 17.

cubes in  $\mathfrak{D}$  contained in  $Q_r$ . Then it is easy to see that  $l - (Q_r^c + \tilde{Q} + Q_j + \dots + Q_s) \approx l_i$ . Therefore  $\tilde{\mathfrak{D}} = \{(l_1, Q_1), \dots, (l - (Q_r^c + \tilde{Q} + Q_j + \dots + Q_s), Q_r), \dots, (l_n, Q_n)\}$ , which is obtained by replacing  $(l_i, \tilde{Q}_i)$  by  $(l - (Q_r^c + \tilde{Q} + Q_j + \dots + Q_s), Q_r)$ , is obviously a decomposition system of  $l$  equivalent to  $\mathfrak{D}$ .

Case (B). By the definition of the decomposition system, one of  $Q_1^c, Q_2^c, \dots, Q_k^c$  must coincide with some cube of  $\mathfrak{D}$ . We may assume without loss of generality that  $Q_1^c = Q_1$ . Then it is easy to see that  $\tilde{l} \approx l_1$ . Replacing  $(l_1, Q_1)$  by  $(\tilde{l}, \tilde{Q})$  and  $(l_i, Q_i)$  by  $(l - (\tilde{Q}^c + Q_j + \dots + Q_s), \tilde{Q}^c)$ , we obtain a decomposition system  $\tilde{\mathfrak{D}}$  of  $l$  which is equivalent to  $\mathfrak{D}$  (Fig. 17). Thus in Case II (1), the cut line  $c$  is deleted without getting new cut lines on  $\dot{Q}_1'$ .

Case II (2).  $E \subset Q_i^c$ .  $A, B \in E$ .

We need not consider this case. For, if we choose any other one of the innermost cut lines on  $\dot{Q}_1'$ , then we can delete it, because it belongs to Case I (0) or Case II (0). Next we delete another one of the innermost cut lines. Thus finally there remains only one cut line  $c$  on  $\dot{Q}_1'$ . The  $c$  is a cut line of Case I (0). Hence we can delete  $c$  from  $\dot{Q}_1'$ .

Repeating step by step the above deleting process, we can delete finally all the cut lines from  $\dot{Q}_1'$ . Let the decomposition system of  $l$  finally obtained be  $\mathfrak{D}'' = \{(l_1'', Q_1''), \dots, (l_n'', Q_n'')\}$ . Then  $\dot{Q}_1' \cap (\bigcup_i \dot{Q}_i) = \emptyset$ . From the condition (D V) it follows directly that there is at least one cube of  $\mathfrak{D}$ , say  $Q_1''$ , such that  $Q_1'' \cap Q_1' \neq \emptyset$ . Then the following two cases are to be considered:

(1)  $Q_1'' \subset \text{Int } Q_1'$ . (2)  $Q_1' \subset \text{Int } Q_1''$ .

Case (1). Let  $l \cap Q_1''$  represents  $\tilde{l}_1$ . Then clearly  $l_1' = \tilde{l}_1 \cdot (l - Q_1'' + Q_1'')$ . Since  $l_1'$  is prime and  $\tilde{l}_1$  is non trivial,  $l - (Q_1'' + Q_1'')$  must be trivial, accordingly  $l_1' = \tilde{l}_1$  as shown in Fig. 18. Then there cannot exist any cube of  $\mathfrak{D}''$  in  $\text{Int } Q_1''$ . Hence  $l \approx \tilde{l}_1''$  and  $(l_1'', Q_1'')$  is an end of  $\mathfrak{D}''$  and obviously  $l_1'' = l \cap Q_2'' = l \cap Q_1' = l_1'^*$ . Thus in Case (1), Lemma 3 is proved.

Case (2). We may suppose that there is no  $Q_1''$  of  $\mathfrak{D}''$  such that  $Q_1' \subset \text{Int } Q_1''$ . If  $(l_1'', Q_1'')$  is an end of  $\mathfrak{D}''$ , then it is easy to see that  $l_1' \approx l_1''$ , since the circumstances are the same as Case (1). Therefore  $(l_1'', Q_1'') \approx (l_1', Q_1')$  and  $l_1'^* \subset l_1''^*$ .

Hence suppose that  $(l_1'', Q_1'')$  is not an end of  $\mathfrak{D}''$ . Let  $Q_1'', \dots, Q_j''$  be all the cubes of  $\mathfrak{D}''$  contained in  $Q_1''$ . Then take out one of the maximal cubes from  $Q_1'', \dots, Q_j''$  and put it  $Q_k''$ , i.e. there is no  $Q_j''$  such that  $Q_k'' \subset Q_j'' \subset Q_1''$ . Then it is easy to see that  $l_1'' \approx l_1'$  and  $l_1'' \approx l - (Q_1'' + Q_1' + \dots + Q_1' + \dots + Q_j'')$  as shown in Fig. 19. Therefore

replacing  $(l_k'', Q_k'')$  by  $(l - (Q_1''^c + Q_i'' + \dots + Q_1' + \dots + Q_j''), Q_1'')$  and  $(l_1', Q_1')$  by  $(l_1', Q_1')$  in  $\mathfrak{D}''$ , we have a decomposition system of  $l$  equivalent to  $\mathfrak{D}''$ , which we write still by  $\mathfrak{D}''$ . Then  $(l_1', Q_1')$  is an end of  $\mathfrak{D}''$ . Hence Lemma 3 is also proved in Case (2). Thus the proof of Lemm 3 is complete.

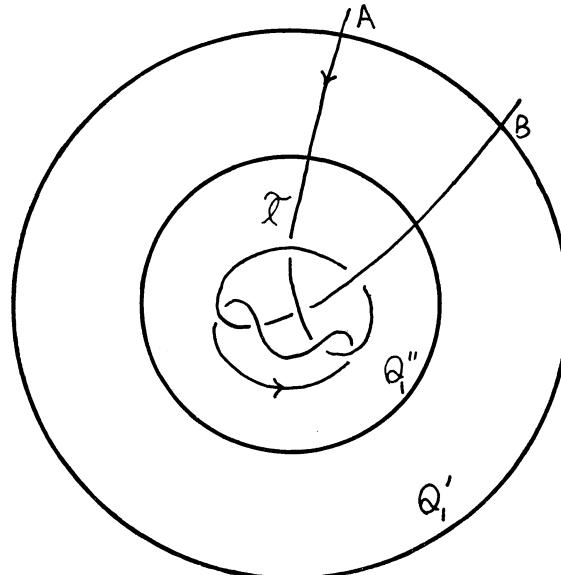


Fig. 18.

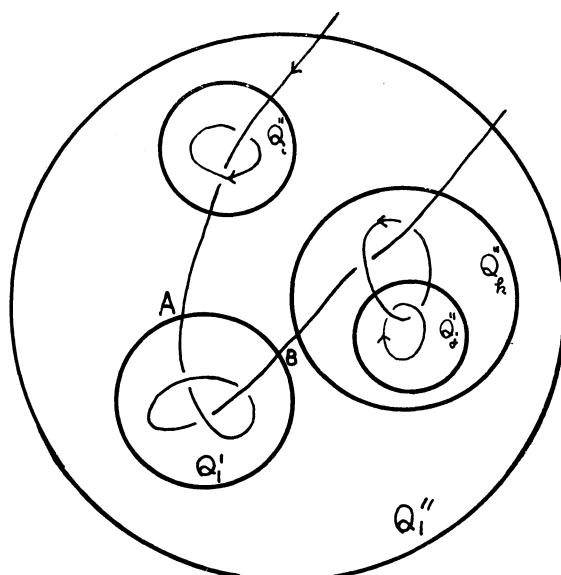


Fig. 19.

**Main Theorem.** *Every non trivial and non separable link can be decomposed uniquely into prime links.*

Proof. Let  $l$  be an arbitrary given non trivial and non separable link and let  $\mathfrak{D} = \{(l_1, Q_1), \dots, (l_n, Q_n)\}$ ,  $\mathfrak{D}' = \{(l'_1, Q'_1), \dots, (l'_m, Q'_m)\}$  be decomposition systems of  $l$ . Then by the use of Lemma 3 we can find an end of  $\mathfrak{D}$ , say  $(l_1, Q_1)$ , and an end of  $\mathfrak{D}'$ , say  $(l'_1, Q'_1)$ , such that  $(l_1, Q_1) \approx (l'_1, Q'_1)$  and one of  $l_1^*$ ,  $l'_1*$  is contained in the other. It is obviously seen that  $l - Q_1 \approx l - Q'_1$  and  $l \approx (l - Q_1) \cdot l_1 \approx (l - Q'_1) \cdot l'_1$ . Then  $\tilde{\mathfrak{D}} = \{(l_2, Q_2), \dots, (l_n, Q_n)\}$  and  $\tilde{\mathfrak{D}'} = \{(l'_2, Q'_2), \dots, (l'_m, Q'_m)\}$  are decomposition systems of  $l - Q_1$  and  $l - Q'_1$  respectively. As is well known, there exists a semilinear mapping  $\varphi$  which maps  $l - Q'_1$  onto  $l - Q_1$ . Then the image  $\varphi\tilde{\mathfrak{D}'} = \{(\varphi l'_2, \varphi Q'_2), \dots, (\varphi l'_m, \varphi Q'_m)\}$  of  $\tilde{\mathfrak{D}'}$  by  $\varphi$  is also a decomposition system of  $l - Q_1$  and clearly  $\varphi\tilde{\mathfrak{D}'} \approx \tilde{\mathfrak{D}'}$ . Next we apply the above method to  $\tilde{\mathfrak{D}}$  and  $\varphi\tilde{\mathfrak{D}'}$ . Thus taking out the equivalent ends from  $\mathfrak{D}$  and  $\mathfrak{D}'$  step by step, we obtain finally the conclusion of the Main Theorem.

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