

Title	Gentzen method in modal calculi. II
Author(s)	Ohnishi, Masao; Matsumoto, Kazuo
Citation	Osaka Mathematical Journal. 1959, 11(2), p. 115-120
Version Type	VoR
URL	https://doi.org/10.18910/5618
rights	
Note	

Osaka University Knowledge Archive : OUKA

<https://ir.library.osaka-u.ac.jp/>

Osaka University

Gentzen Method in Modal Calculi, II¹⁾

To Professor Zyoiti SUETUNA on his 60th birthday

By Masao OHNISHI and Kazuo MATSUMOTO

Various decision procedures for modal sentential calculus S5 have been given by W. T. Parry [8], R. Carnap [10], M. Itoh [11], M. Ohnishi-K. Matsumoto [7] and S. Kanger [5]. Among them Gentzen-type procedure is only that of S. Kanger. The object of this note is to give an alternative Gentzen-type decision procedure for S5.

I

Our formulation of Lewis's modal sentential calculus S5²⁾ is based upon "Sequenzenkalkül *LK*", which was constructed by G. Gentzen [3]. Namely :

$\left\{ \begin{array}{l} \text{logical symbols :} \\ \quad \cdot \text{ (and), } \sim \text{ (not), } \vee \text{ (or)} \\ \text{rules of inference :} \\ \quad \left\{ \begin{array}{l} \text{structural rules} \\ \quad \text{weakening, contraction, exchange and cut.} \\ \text{logical rules} \\ \quad (\rightarrow \cdot), (\cdot \rightarrow); (\rightarrow \sim), (\sim \rightarrow); (\rightarrow \vee), (\vee \rightarrow). \end{array} \right. \end{array} \right.$

Next we add to *LK* a new logical symbol \Box (necessary), and we define as follows: if α is a formula, then $\Box\alpha$ is also a formula.

New rules for modality are

$$\frac{\alpha, \Gamma \rightarrow \Theta}{\Box\alpha, \Gamma \rightarrow \Theta} (\Box \rightarrow), \quad \frac{\Box\Gamma \rightarrow \Box\Theta, \alpha}{\Box\Gamma \rightarrow \Box\Theta, \Box\alpha} (\rightarrow \Box).$$

By Γ, Θ we mean a series of formulas as in *LK*. $\Box\Gamma$ (or $\sim\Gamma$) means a series of formulas which is formed by prefixing \Box (or \sim) in front of each formula of Γ .

1) This is a continuation to M. Ohnishi and K. Matsumoto [7].

2) C. I. Lewis and C. H. Langford [6].

Thus established sentential calculus which contains LK is $S5^3$.

M. Wajsberg [9] gave a decision procedure for monadic functional calculus, and W. T. Parry [8] remarked that for an arbitrary formula γ in $S5$ there exists a γ^* equivalent to γ and of degree at most 1⁴, and using this fact he showed the equivalence of $S5$ and monadic functional calculus, hence gave a decision procedure for $S5$.

In the following γ^* always means a formula of degree at most 1 and equivalent to γ .

When all sequent-formulas of each rule in $S5$ are of degree at most 1, we denote this system by $S5^*$.

Theorem 1. $\rightarrow\gamma$ is provable in $S5$, if and only if $\rightarrow\gamma^*$ is provable in $S5^*$.

Lemma 1. Let $\rightarrow\gamma$ be an $S5$ -provable sequent. Then there exists an $S5$ -proof-figure for $\rightarrow\gamma$ which does not include any mix other than \Box -mix⁵.

The proof of Lemma 1 is carried out by the induction on the rank and grade of mix-formula.

We define a modalized formula (abbr. mf) as follows: (1-3)

1. $\Box\alpha$ is an mf.
2. If α is an mf, then so is $\sim\alpha$.
3. If α and β are mfs, then so are $\alpha\vee\beta$ and $\alpha\cdot\beta$.

Clearly we have

Lemma 2⁶. In the formulation of $S5$ we can replace $(\rightarrow\Box)$ by

$$\frac{\Gamma' \rightarrow \Theta', \alpha}{\Gamma' \rightarrow \Theta', \Box\alpha} (\rightarrow\Box)',$$

where Γ' and Θ' mean series of mfs.

Lemma 3. Assuming that all mixes appearing in an $S5$ -proof-figure of $\rightarrow\gamma$ are \Box -mixes, there exists a proof-figure of $\rightarrow\gamma$, where the degrees of mix-formulas are all at most 1.

3) The following example shows that Gentzen's Hauptsatz does not hold in our $S5$:

$$\frac{\frac{\frac{\Box\alpha \rightarrow \Box\alpha}{\rightarrow\Box\alpha, \sim\Box\alpha} (\rightarrow\sim) \quad \frac{\alpha \rightarrow \alpha}{\Box\alpha \rightarrow \alpha} (\Box\rightarrow)}{\rightarrow\Box\alpha, \Box\sim\Box\alpha} (\rightarrow\Box) \quad \frac{\alpha \rightarrow \alpha}{\Box\alpha \rightarrow \alpha} (\Box\rightarrow)}{\rightarrow\Box\sim\Box\alpha, \alpha} (\Box\alpha)$$

4) For the definition of "a formula of degree n ", see A. R. Anderson [1], p. 203.
 5) \Box -mix means a mix, the outermost symbol of whose mix-formula is \Box .
 6) This Lemma is a formal generalization of R. Feys's formulation [2].

The proof of Lemma 3 can be carried out by the induction on n , where n is the maximal degree of \Box -mix-formula appearing in the proof-figure of $\rightarrow\gamma$. Let $\Box\xi$ be an upper-most (in the proof-figure) mix-formula of degree n (2). By the aid of the following Wajsberg's recurring equivalences⁷⁾: $\Box\Box\alpha = \Box\alpha$, $\Box\sim\Box\alpha = \sim\Box\alpha$, $\Box(\Box\alpha \vee \beta) = \Box\alpha \vee \Box\beta$, $\Box(\sim\Box\alpha \vee \beta) = \sim\Box\alpha \vee \Box\beta$ and $\Box(\alpha \cdot \beta) = \Box\alpha \cdot \Box\beta$, clearly we can define ξ^\dagger with the following properties: 1. ξ^\dagger is an mf equivalent to $\Box\xi$, 2. ξ^\dagger is of degree at most $n-1$, and 3. in the upper part of the \Box -mix of $\Box\xi$ we can replace $\Box\xi$ by ξ^\dagger without any \Box -mix of degree n .

Now, Theorem 1 can be proved by Lemmas 1, 2 and 3.

Theorem 2. *The Hauptsatz does hold in S5*.*

In order to prove Theorem 2 we need the following

Lemma 4. *In S5*, if $\Box\Gamma \rightarrow \Box^\Theta, \alpha$ is provable without any mix, then either $\Box\Gamma \rightarrow \Box^\Theta$ or $\Box\Gamma \rightarrow \alpha$ is provable without any mix, where all formulas of Γ , Θ and α are of degree 0 (i.e. LK-formulas).*

The proof of Lemma 4 can be easily given by the idea of elimination of formula-bundle of each formula of \Box^Θ in the proof-figure of $\Box\Gamma \rightarrow \Box^\Theta, \alpha$.

Now we have only to consider the following cases: 1. When $\rho=2$ and the mix is a \Box -mix, i.e.

$$\frac{\frac{\Box\Gamma \rightarrow \Box^\Theta, \alpha}{\Box\Gamma \rightarrow \Box^\Theta, \Box\alpha} (\rightarrow \Box) \quad \frac{\alpha, \Sigma \rightarrow \Pi}{\Box\alpha, \Sigma \rightarrow \Pi} (\Box \rightarrow)}{\Box\Gamma, \Sigma \rightarrow \Box^\Theta, \Pi} (\Box\alpha),$$

we transform this into:

$$\frac{\frac{\Box\Gamma \rightarrow \Box^\Theta, \alpha}{\Box\Gamma, \Sigma^* \rightarrow (\Box^\Theta)^*, \Pi} (\alpha)}{\Box\Gamma, \Sigma \rightarrow \Box^\Theta, \Pi}$$

2. When $\rho > 2$ and the left rank > 1 ,

$$2.1 \quad \frac{\frac{\alpha, \Gamma \rightarrow \Theta}{\Box\alpha, \Gamma \rightarrow \Theta} (\Box \rightarrow) \quad \Sigma \rightarrow \Pi}{\Box\alpha, \Gamma, \Sigma^* \rightarrow \Theta^*, \Pi} (\mathfrak{M}).$$

This case is trivial.

$$2.2 \quad \frac{\frac{\Box\Gamma \rightarrow \Box^\Theta, \alpha}{\Box\Gamma \rightarrow \Box^\Theta, \Box\alpha} (\rightarrow \Box) \quad \Sigma \rightarrow \Pi}{\Box\Gamma, \Sigma^* \rightarrow (\Box^\Theta)^*, (\Box\alpha)^*, \Pi} (\Box\mathfrak{M}).$$

7) See M. Wajsberg [9].

According to Lemma 4 either $\Box\Gamma \rightarrow \Box\Theta$ or $\Box\Gamma \rightarrow \alpha$ is provable. In the former case 2.21 we transform this into:

$$\frac{\frac{\Box\Gamma \rightarrow \Box\Theta \quad \Sigma \rightarrow \Pi}{\Box\Gamma, \Sigma^* \rightarrow (\Box\Theta)^*, \Pi} (\Box\mathfrak{M})}{\Box\Gamma, \Sigma^* \rightarrow (\Box\Theta)^*, (\Box\alpha)^*, \Pi} \text{ (weakening if necessary).}$$

In the latter case 2.22 we distinguish the following two cases:

2.221. when $\mathfrak{M} = \alpha$,

$$\frac{\frac{\frac{\Box\Gamma \rightarrow \alpha}{\Box\Gamma \rightarrow \Box\alpha} (\rightarrow \Box) \quad \Sigma \rightarrow \Pi}{\Box\Gamma, \Sigma^* \rightarrow \Pi} (\Box\alpha),}{\Box\Gamma, \Sigma^* \rightarrow (\Box\Theta)^*, \Pi}$$

2.222. when $\mathfrak{M} \neq \alpha$,

$$\frac{\frac{\frac{\Box\Gamma \rightarrow \alpha}{\Box\Gamma \rightarrow \Box\alpha} (\rightarrow \Box)}{\Box\Gamma \rightarrow \Box\Theta, \Box\alpha} \quad \Sigma \rightarrow \Pi}{\Box\Gamma, \Sigma^* \rightarrow (\Box\Theta)^*, \Box\alpha, \Pi} (\Box\mathfrak{M}).$$

3. When $\rho > 2$, the left rank = 1 and the right rank > 1 ,

3.1

$$\frac{\Sigma \rightarrow \Pi \quad \frac{\alpha, \Gamma \rightarrow \Theta}{\Box\alpha, \Gamma \rightarrow \Theta} (\Box \rightarrow)}{\Sigma, (\Box\alpha)^*, \Gamma^* \rightarrow \Pi^*, \Theta} (\mathfrak{M}),$$

we transform this into: 3.11. in case $\mathfrak{M} = \Box\alpha$,

$$\frac{\Sigma \rightarrow \Pi \quad \frac{\frac{\Sigma \rightarrow \Pi \quad \alpha, \Gamma \rightarrow \Theta}{\Sigma, \alpha, \Gamma^* \rightarrow \Pi^*, \Theta} (\Box\alpha)}{\Sigma, \Box\alpha, \Gamma^* \rightarrow \Pi^*, \Theta} (\Box \rightarrow)}{\Sigma, \Sigma^*, \Gamma^* \rightarrow \Pi^*, \Pi^*, \Theta} (\Box\alpha)}{\Sigma, \Gamma^* \rightarrow \Pi^*, \Theta}$$

3.12. in case $\mathfrak{M} \neq \Box\alpha$,

$$\frac{\frac{\Sigma \rightarrow \Pi \quad \alpha, \Gamma \rightarrow \Theta}{\Sigma, \alpha^*, \Gamma^* \rightarrow \Pi^*, \Theta} (\mathfrak{M})}{\Sigma, \alpha, \Gamma^* \rightarrow \Pi^*, \Theta} \text{ (weakening if necessary)} \\ \frac{\Sigma, \alpha, \Gamma^* \rightarrow \Pi^*, \Theta}{\Sigma, \Box\alpha, \Gamma^* \rightarrow \Pi^*, \Theta} (\Box \rightarrow)$$

3.2

$$\frac{\Sigma \rightarrow \Pi \quad \frac{\Box\Gamma \rightarrow \Box\Theta, \alpha}{\Box\Gamma \rightarrow \Box\Theta, \Box\alpha} (\rightarrow \Box)}{\Sigma, (\Box\Gamma)^* \rightarrow \Pi^*, \Box\Theta, \Box\alpha} (\Box\mathfrak{M}),$$

where Π contains $\Box\mathfrak{M}$, and the left rank = 1. Therefore the non-trivial case is

$$\frac{\frac{\frac{\Box\Sigma \rightarrow \Box\Pi, \mathfrak{M}}{\Box\Sigma \rightarrow \Box\Pi, \Box\mathfrak{M}} (\rightarrow \Box) \quad \frac{\Box\Gamma \rightarrow \Box\Theta, \alpha}{\Box\Gamma \rightarrow \Box\Theta, \Box\alpha} (\rightarrow \Box)}{\Box\Sigma, (\Box\Gamma)^* \rightarrow \Box\Pi, \Box\Theta, \Box\alpha} (\Box\mathfrak{M}) .$$

e transform this into :

$$\frac{\frac{\Box\Sigma \rightarrow \Box\Pi, \Box\mathfrak{M} \quad \Box\Gamma \rightarrow \Box\Theta, \alpha}{\Box\Sigma, (\Box\Gamma)^* \rightarrow \Box\Pi, \Box\Theta, \alpha} (\Box\mathfrak{M})}{\Box\Sigma, (\Box\Gamma)^* \rightarrow \Box\Pi, \Box\Theta, \Box\alpha} (\rightarrow \Box) .$$

This completes the proof of Theorem 2.

Theorems 1 and 2 yield a new decision procedure for S5.

II

W. T. Parry [8] showed that if a formula γ of degree 1 is provable in S5, then γ is also provable in S3. S. Halldén [4] remarked that if a formula γ of degree 1 is provable in any one system of S2, S3, S4 and S5, then γ is also provable in each of other systems. That is,

Theorem 3. $\rightarrow\gamma^*$ is provable in S5, if and only if $\rightarrow\gamma^*$ is provable in S2, where γ^* is of degree 1.

In our former paper [7], we have already obtained the following result: “ $\rightarrow\gamma$ is provable in S2 if and only if $p\text{-}\exists p\text{-}\rightarrow\gamma^8$ is provable in Q2, where p is a sentence-variable”.

This result leads to an alternative proof of Theorem 3. That is, we have only to show that $p\text{-}\exists p\text{-}\rightarrow\gamma^*$ is provable in Q2 if $\rightarrow\gamma^*$ is provable in S5* (see Theorem 1). The essential part of this proof is to show the Q2-admissibility of

$$\frac{p\text{-}\exists p, \Box\Gamma \rightarrow \Box\Theta, \alpha}{p\text{-}\exists p, \Box\Gamma \rightarrow \Box\Theta, \Box\alpha} ,$$

where Γ, Θ and α are all LK-formulas. According to an analogous form of Lemma 4 if $p\text{-}\exists p, \Box\Gamma \rightarrow \Box\Theta, \alpha$ is provable in Q2, then either $p\text{-}\exists p, \Box\Gamma \rightarrow \Box\Theta$ or $p\text{-}\exists p, \Box\Gamma \rightarrow \alpha$ is provable in Q2. The former case is clear. In the latter case eliminating all \Box 's in the proof of $p\text{-}\exists p, \Box\Gamma \rightarrow \alpha$, we obtain a proof of $p\text{-}\exists p, \Gamma \rightarrow \alpha$, hence of $p\text{-}\exists p, \Box\Gamma \rightarrow \Box\alpha$.

(Received July 27, 1959)

8) $\alpha\text{-}\exists\beta$ is the abbreviation of $\Box(\sim\alpha \vee \beta)$.

Bibliography

- [1] A. R. Anderson: Improved decision procedures for Lewis's calculus S4 von Wright's calculus M. J. Symbolic Logic **19**, 201-214 (1954).
- [2] R. Feys: Les systèmes formalisée des modalités aristotéliennes, Revue Phil. de Louvain **48**, 478-509 (1950).
- [3] G. Gentzen: Untersuchungen über das logische Schliessen, I. II., Math. Z. **39**, 176-210, 405-431 (1935).
- [4] S. Halldén: Results concerning the decision problem of Lewis's calculi S3 and S6, J. Symbolic Logic **14**, 230-236 (1950).
- [5] S. Kanger: Provability in Logic, Stockholm (1957).
- [6] C. I. Lewis and C. H. Langford: Symbolic Logic, New York (1932).
- [7] M. Ohnishi and K. Matsumoto: Gentzen Method in Modal Calculi, Osaka Math. J. **9**, 113-130 (1957).
- [8] W. T. Parry: Zum Lewisschen Aussagenkalkül, Erg. eines math. Koll. **4**, 15 (1931-2).
- [9] M. Wajsberg: Ein erweiterter Klassenkalkül, Monats. f. math. u. phys. **40**, 113-126 (1933).
- [10] R. Carnap: Modalities and Quantification, J. Symbolic Logic, **11**, 33-64 (1946).
- [11] M. Itoh: On the relation between the modal sentential logic and the monadic predicate calculus, I. II. III. J. Japan Ass. Phil. of Sci. **3**, 142-145, **4**, 162-167 (1955), **6**, 258-265 (1956) (in Japanese).